## CS 6160: Voronoi Diagrams

Due Date: .

This assignment has 6 questions, for a total of 100 points and 20 bonus points. Unless otherwise specified, complete and reasoned arguments will be expected for all answers.

| Question | Points | Bonus Points | Score |
| :---: | :---: | :---: | :---: |
| Distances | 20 | 0 |  |
| Properties | 30 | 0 |  |
| Reconstruction | 20 | 0 |  |
| Farthest Point Diagrams | 20 | 0 |  |
| Minimum Enclosing ball | 10 | 0 |  |
| More Reconstruction | 0 | 20 |  |
| Total: | 100 | 20 |  |

Question 1: Distances
For each of the following distance functions, describe what the corresponding Voronoi diagram will look like for $n$ points in the plane. Your answer should describe what the 0 and 1 dimensional parts of the diagram look like, as well as give an estimate of the overall complexity of the diagram.

- $d(\mathbf{p}, \mathbf{p})=\left|p_{x}-q_{x}\right|+\left|p_{y}-q_{y}\right|$
- $d(\mathbf{p}, \mathbf{p})=\max \left(\left|p_{x}-q_{x}\right|,\left|p_{y}-q_{y}\right|\right)$

Question 2: Properties.
(a) [10] Prove that a Voronoi cell is unbounded if and only if the site associated with that cell lies on the convex hull of the set of sites.
(b) [10] Prove that a Voronoi cell under the Euclidean distance must be convex.
(c) [5] Can you describe a distance function for which a Voronoi cell is not convex?
(d) [5] Argue for why computing the Voronoi diagram in the plane must take at least $\Omega(n \log n)$ time.

Question 3: Reconstruction
Any convex polygon can be realized as a Voronoi cell in the Voronoi diagram of a set of points. Prove this fact by taking any convex polygon and constructing a set of points for which this polygon is one Voronoi cell.

Question 4: Farthest Point Diagrams
Design an efficient algorithm to compute a farthest point Voronoi diagram. Note that only points on the convex hull will have cells in the diagram. You might consider a divide-and-conquer strategy.

Question 5: Minimum Enclosing ball.
Let the minimum enclosing ball of a set of points $P$ be a ball of radius $r$ that encloses all of $P$ and where $r$ is as small as possible. Assuming you're given an algorithm that can compute the furthest-point Voronoi diagram in $O(n \log n)$ time. Present an algorithm that computes the minimum enclosing ball in $O(n \log n)$ time. You can assume that the Voronoi diagram is appropriately labeled so that for each cell you know which site its points are furthest from.

BONUS Question 6: More Reconstruction
Suppose you are given a Voronoi diagram in which each vertex has degree exactly three (i.e a nondegenerate case). However, you are not given the actual $n$ sites from which the diagram was constructed. Design an algorithm that runs in linear time to reconstruct a set of sites consistent with this diagram.
HINT: Consider a single cell and a Voronoi vertex on its boundary. Find a way to determine the site corresponding to that cell.

