CS 5110/6110 – Rigorous System Design | Spring 2016 Feb-23

#### Lecture 14 Concolic Execution

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#### Last Time

- Drawbacks of concrete testing
- Symbolic execution
- Solutions for the path explosion problem
  - Structural abstraction
  - Compositional symbolic execution

## Symbolic Execution

- Key idea: execution of programs using symbolic input values instead of concrete data
- Concrete vs symbolic
  - Concrete execution
    - Program takes only one path determined by input values
  - Symbolic execution
    - Program can take any feasible path coverage!

# Symbolic Program State

- Symbolic values of program variables
- Path condition (PC)
  - Logical formula over symbolic inputs
  - Accumulates constraints that inputs have to satisfy for the particular path to be executed
  - If a path is feasible its PC is satisfiable
- Program location

## Symbolic Execution Tree

- Characterizes execution paths constructed during symbolic execution
- Nodes are symbolic program states
- Edges are labeled with program transitions

#### Example

1) **int** x, y; 2) if (x > y) {  $\mathbf{x} = \mathbf{x} + \mathbf{y};$ 3) 4) y = x - y;5) X = X - Y;6) **if** (x > y)assert false; 7) 8) }



#### Further Limitations of Symbolic Execution

- Limited by the power of constraint solver
  - Cannot handle non-linear and very complex constraints
- Inherently white-box technique
  - Source code (or equivalent) is required for precise symbolic execution
  - Modeling libraries is a huge problem

## This Time

- Combining concrete and symbolic execution
- Many names referring to the same thing:
  - DART (directed automated random testing)
  - Concolic (concrete + symbolic) execution
  - Dynamic symbolic execution

#### **Concolic Execution**

- Combination of concrete and symbolic execution to overcome the two weaknesses of classic symbolic execution
- Algorithm
  - Execute program concretely
  - Collect the symbolic path condition along the way
  - Negate a constraint on the path condition after the run and solve it to get a model
  - Execute again with the newly found concrete input values

```
Simple Example I
void foo(int x, int y) {
  z = 2*y;
  if (z == x) {
    if (x > y+10) {
      assert false;
    }
  }
```

Simple Example II

```
void foo(int x, int y) {
  z = double(y); // no source or complex
  if (z == x) {
    if (x > y+10) {
      assert false;
    }
```

## **High-Level Picture**



## **High-Level Picture**



## **High-Level Picture**



## **Concolic Covering Middle Ground**

#### Concrete

- + Complex programs
- + Binaries
- + Scalable
- Less coverage

+ No false positives

+ Complex programs
+ Binaries
+/- Scalable
+ High coverage

Concolic

+ No false positives

#### Symbolic

- Simple programs
- Source code
- Not scalable
- + High coverage
- False positives

#### **Recent Success Stories**

#### SAGE

- Microsoft's internal tool for finding security bugs
- White-box fuzzing
  - Concolic execution for finding bugs in file parsers (jpeg, docx, ppt,...)
- Last line of defense
- Big clusters continuously running SAGE

#### KLEE

- Open source concolic executor
- Runs on top of LLVM
- Has found lots of problems in open-source software